

Tutorial LTL CTL

Question 1

- ① $G F m$ {liveliness} model satisfies
- ② $F G c$ {safety} model fails
- ③ $G [c \rightarrow X [F m \wedge (\neg m \cup g)]]$ {safety} model fails

Question 2

- ① $E F A G \bar{a}$ [holds on none of the states]
- ② $E G \bar{E}$ [holds on m and h]
- ③ $E (\bar{a} \cup E)$ [holds on h, c, g]
- ④ $A G (a \vee E G (\neg E))$ [does not hold on any state]
- ⑤ $E G [a \vee E G (\neg E)]$ [only c does not satisfy]
- ⑥ $A G A F (a \vee \neg E)$ [holds on all paths]

Question 3

① $A \wedge G [A \rightarrow X B \wedge X X B] \quad \Sigma = \{A, B, AB, \phi\}$
 $\omega\text{-RE} = [AB]^\omega$

② $G (A \rightarrow X ((\neg A \cup (B \wedge \neg A)) \vee G \neg A)) \quad \Sigma = \{A, B, C, AB, AC, BC, ABC, \phi\}$
 $\omega\text{-RE} : (B+C)^* [A \cdot (B+C)^* B (B+C)^* + B+C]^\omega$

③ $G (X A \wedge X X B \rightarrow C)$
 $\omega\text{-RE} : C B + C + A + C + (AB)^\omega + (BA)^\omega$

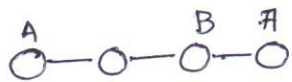
④ $G [\{(\neg A \wedge B) \vee C\} \Rightarrow X F (A \cap B)] \quad AP = \{A, B, C\}$
 $\Sigma = 2^{AP} = \{A, B, C, AB, BC, AC, ABC, \phi\}$
 $\omega\text{-RE} : \{ (A+B)^* + [(A+B)^* B^* + T^* X] \}^\omega \quad d \subseteq \Sigma = \{AB, X\} \text{ where } X = AB$

Question 4 Equivalence checking

$$\textcircled{1} \quad G(A \rightarrow B) \equiv G A \rightarrow G B$$

not equivalent

counter example



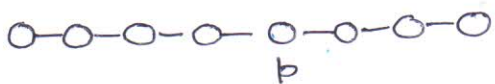
↳ L.H.S fails

R.H.S bases (vacuously true)

$$\textcircled{2} \quad G F p \rightarrow G F q \equiv G(p \rightarrow F q)$$

not equivalent

counter example:



p only occurs once in the trace π and q never occurs

$G F p$ is false so L.H.S is vacuously true

When p occurs $F q$ does not occur so R.H.S is false

$$3. \quad FGp \wedge FGq \equiv FG(p \wedge q)$$

equivalent, Let path $\pi = s_0 s_1 s_2 \dots$

Part 1: LHS \rightarrow R.H.S

$$\pi \vDash FGp \wedge FGq$$

$$\Rightarrow (\exists a \geq 0, \forall i \geq a, s_i \vDash p) \wedge (\exists b \geq 0, \forall j \geq b, s_j \vDash q)$$

Case 1 $a \geq b$

$$\forall i \geq a, s_i \vDash p \ \& \ s_i \vDash q$$

$$\Rightarrow \forall i \geq a, s_i \vDash (p \wedge q)$$

$$\Rightarrow \pi \vDash FG(p \wedge q)$$

Case 2 $b > a$

$$\forall i \geq b, s_i \vDash p \ \& \ s_i \vDash q$$

$$\Rightarrow \forall i \geq b, s_i \vDash (p \wedge q)$$

$$\Rightarrow \pi \vDash FG(p \wedge q)$$

\therefore if $\pi \vDash (FGp \wedge FGq)$, then $\pi \vDash FG(p \wedge q)$ — (1)

Part 2: R.H.S \rightarrow L.H.S

$$\pi \vDash FG(p \wedge q)$$

$$\Rightarrow \exists a \geq 0, \forall i \geq a, s_i \vDash (p \wedge q)$$

$$\Rightarrow \exists a \geq 0, \forall i \geq a, s_i \vDash p \ \& \ s_i \vDash q$$

$$\therefore a \geq 0, s_a \vDash FGp \ \text{and} \ s_a \vDash FGq$$

$$\Rightarrow s_a \vDash (FGp \wedge FGq) \Rightarrow \pi \vDash (FGp \wedge FGq)$$

\therefore if $\pi \vDash FG(p \wedge q)$, then $\pi \vDash (FGp \wedge FGq)$ — (2)

From (1) & (2) $FGp \wedge FGq \equiv FG(p \wedge q)$